## Entropy of Turing machines with moving head

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In Kůrka [1] the last theorem says that a Turing machine with moving head (TMH) has the zero topological entropy, but the proof is not correct. I give here a correct proof. Given a Turing machine with alphabet A, set of inner states Q and the transition function  $\delta = (\delta_Q, \delta_A, \delta_Z) : Q \times A \rightarrow Q \times A \times \{-1, 0, 1\}$ , we have the sofic subshift

$$X_1 = \{ x \in B^{\mathbb{Z}} : |\{ i \in \mathbb{Z} : x_i \in Q \}| \le 1 \}$$

of all configurations which contain at most one head. The TMH determined by  $\delta$  is a cellular automaton  $F: X_1 \to X_1$  (see [1] for the definition of F).

Given an initial configuration  $x \in A^{\mathbb{Z}}$  and an interval  $I \subseteq \mathbb{Z}$  of width |I| = d, we consider rectangles  $I \times [kd^2, (k+1)d^2)$  in the space-time diagram. Since there are at most  $d^2$  heads in such a rectangle, one of the d columns of I contains at most d heads. We say that an I-cut is a collection of these columns, one for each rectangle  $I \times [kd^2, (k+1)d^2)$ . Formally an *I*-cut of x is a sequence of integers  $(j_k \in I)_{k>0}$  such that for each  $k \ge 0$  we have  $|\{i \in [kd^2, (k+1)d^2) : F^i(x)_{j_k} \in Q\}| \le d$ . Clearly, for each x and I, an I-cut exists. Given arbitrary finite interval  $J \subseteq \mathbb{Z}$  we estimate the entropy of the column subshift  $\Sigma_J = \{ u \in (B^{|J|})^{\mathbb{N}} : \exists x \in X, \forall i, u_i = F^i(x)_J \} \subseteq (B^{|J|})^{\mathbb{N}}.$  Given d > 0, let  $I_0, I_1$  be the columns of width d which are left and right neighbours of J (and disjoint with J). A word of  $\mathcal{L}(\Sigma_J)$  of length  $nd^2$  obtained from  $x \in X$  is determined by the top row  $(x_J) \in B^{|J|}$  and by the values of cuts of length  $nd^2$  through intervals  $I_0, I_1$ . Here we define the value of a cut  $j_k$  in the rectangle  $I \times [kd^2, (k+1)d^2)$  as consisting of the value of the top row  $F^{kd^2}(x)_I$  and of the values  $\{F^i(x)_{j_k}: kd^2 \leq i < (k+1)d^2\}$  of the  $j_k$ -th column. The number of values of all possible cuts through  $I \times [kd^2, (k+1)d^2)$  is bounded by the product of the number  $|B|^d$  of values of all possible case through d of possible columns  $j_k$ , number  $(d^2)^d$  of possible d-tuples  $i \in [kd^2, (k+1)d^2)$  with  $F^i(x)_{j_k} \in Q$ , number  $|Q|^d$  of d-tuples the states of the head at these positions, and the number  $|A|^d$  of d-tuples of the scanned letters, i.e., by  $|B|^d \cdot d^{2d+1} \cdot |Q|^d \cdot |A|^d$ . For the number  $P(nd^2)$  of words of  $\mathcal{L}(\Sigma_J)$  of length  $nd^2$  we get

$$\begin{array}{rcl}
P(nd^2) &\leq & |B|^{|J|} \cdot (|B|^d \cdot d^{2d+1} \cdot |Q|^d \cdot |A|^d)^{2n} \\
\frac{\ln P(nd^2)}{nd^2} &\leq & \frac{|J| \cdot \ln |B|}{nd^2} + \frac{2\ln(|B| \cdot |Q| \cdot |A|)}{d} + \frac{2(2d+1)\ln d}{d^2} \\
H_{top}(\Sigma_J) &\leq & \frac{2\ln(|B| \cdot |Q| \cdot |A|)}{d} + \frac{2(2d+1)\ln d}{d^2}
\end{array}$$

Since this holds for arbitrary d > 0, we get  $H_{top}(\Sigma_J) = 0$  for each finite interval  $J \subseteq \mathbb{Z}$  and it follows that the topological entropy of TMH  $F : X_1 \to X_1$  is zero.

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## References

 Petr Kůrka: On topological dynamics of Turing machines. Theoretical Computer Science 174 (1997) 203-216